

Complexity of rational solutions to polynomial equations.

• Rational ~~points~~ solutions are hard.

Examples: • 'Fermat': $x^n + y^n = z^n$ for $n \geq 3$ has ~~only~~ only trivial solns ($xyz=0$) - Taylor, Wiles, '95.

• The (Y given by ^(k3) $1 + zx^4 = y^4 + 4z^4$, known solns:

• $(0, \pm 1, 0)$

• $(\frac{+12031020}{1484801}, \frac{+1169407}{1484801}, \frac{+1157520}{148401})$ (Eisenhans & Zehnel, '05)

• Are there any others? Unknown.

• Does Eqn of this form (~~or~~ $a + bx^4 = cy^4 + dz^4$) with finite non 0 # of solns? Unknown.

• Fibonacci: $F_0 = 0, F_1 = 1, F_{n+2} = F_n + F_{n+1}$.

Thm: Only perfect powers are:

$$F_0 = 0, F_1 = 1, F_2 = 1, F_6 = 8 \text{ \& } F_{12} = 144$$

(Bugeaud, Mignotte, Siksek, '06)

- used modularity (à la Wiles), + class field stuff

(Baker, lin forms, logs)

From now on, Poincaré

Elliptic Curves

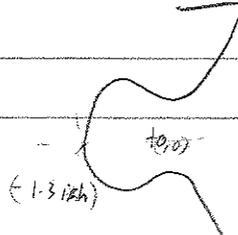
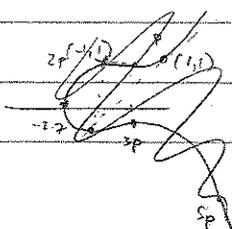
Fix $a, b \in \mathbb{Q}$ st. $4a^3 + 27b^2 \neq 0$ (else ~~same~~ but different). Then the eqn

$$y^2 = x^3 + ax + b \quad \text{is called an elliptic curve}$$

eg: $a = -1, b = 1, y^2 = x^3 - x + 1$ (real root -2.7154)

Draw sketch of real pts - looks like a cune

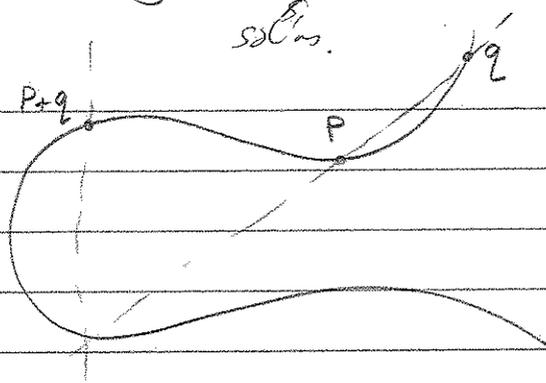
- not like an ellipse (cf. arc lengths)



Main tool for understanding rat. pts on EC:

\uparrow that ∞ .

Group law,



(tangent for $2P$)

This makes {rat. solns to ~~E~~ E } \cup $\{\infty\}$ into an abelian grp, $E(\mathbb{Q})$

Thm [Mordell, Weil, 1923]: $E(\mathbb{Q})$ is a finitely generated.

(so $E(\mathbb{Q}) \cong T \oplus \mathbb{Z}^r$, $\#T$ torsion
 $r = \text{rank}, \in \mathbb{Z}_{\geq 0}$)

- Was hard then, now in 1st EC course for master's students;
- Analogue for K3 surfaces ($\dim CH^2(X)_{\mathbb{Q}} < \infty$) unknown.
(even = 1)

(for $y^2 = x^3 + 1$, $E(\mathbb{Q}) \cong \mathbb{Z}$).

How do T & r vary as we vary a & b ?
torsion rank

• Thm (Nagura 1977) $\#T \leq 12$, $\#$ (sharp)

• Analogue for abelian varieties (higher dim. generalizations of ECs).
still open - 'strong torsion conjecture'

• Unknown if r bounded - ~~was~~ not even a decent conjecture

• Algorithm to compute T .

• BSD \Rightarrow Algorithm for r .

Arithmetic complexity of sol'ns.

• Back to $1+2x^4 = y^4 + 4z^4$: the sol'ns $(0, \pm 1, 0)$ are much 'simpler' than

'Size' not good, $\frac{1}{2}$ vs $\frac{1000}{2001}$. $(\frac{\pm 1203120}{1486201}, \dots)$

• To make precise:

Def: Given $x \in \mathbb{Q}^*$, write $x = \frac{a}{b}$ with a, b coprime, & define

$$h(p) = \log_{\max}(|a|, |b|)$$

Set $h(0) = 0$.

On elliptic curves, extra useful. Given $p = (x_p, y_p) \in E(\mathbb{Q})$, define $h(p) = h(x_p)$.

Eg on $y^2 = x^3 - x + 1$,

$p = (-1, 1)$, so $h(p) = 0$

$p = (3, -5)$, so $h(p) = \log 5$

$2p = (\frac{19}{25}, \frac{103}{125})$, so $h(2p) = \log 25 = 2 \log 5$

$4p = (\frac{-350701}{265225}, \frac{-13919607}{136590875})$, $h(4p) = \log(350701) = 3.97 \cdot h(2p)$

$8p = \dots$, $h(8p) = 4.03 \cdot h(4p)$

In general, $h(2p) \approx 4h(p)$. *Edge*

Thm [Tate]: Let E an EC. Then \exists a constant c & a quadratic form \hat{h} on E such that $\forall p \in E(\mathbb{Q})$, have

$$|h(p) - \hat{h}(p)| \leq c$$

This \hat{h} , the NTht; is

- crucial in proving MW
- essential to formulation of BSD
- non-degenerate on $E(\mathbb{Q}) \otimes \mathbb{Q}$?

Families of varieties

Eg $y^2 = x^3 - t^2x + t^2$. For fixed value of $t \in \mathbb{Q} \setminus \{0\}$, get an EC.
($t = \pm 1$: get $y^2 = x^3 - x + 1$ again).

That EC has \hat{h}_t : how does it vary as we vary t ?
(Fix $t \rightarrow E_t, \hat{h}_t: E_t(\mathbb{Q}) \rightarrow \mathbb{R}$)

[Sample question: how does constant in last term vary with t ?]
• Not v. interesting, some answers, not useful.
• pt is that h is not interesting.

To ask better question, need notion of:

Locally decomposable height.

Def: let X be the set of rat subset of a collection of polynomials.
- eg. $E(\mathbb{Q})$, or \mathbb{Q} itself, or ...
A height h on X is just a fctn $X \rightarrow \mathbb{R}$.

eg. $X = E(\mathbb{Q})$, heights h, \hat{h} , constant fctn to \mathbb{R} , ...
- too many!

Special class of ht functions: locally decomposable hts

~~Def~~ Eg $h: \mathbb{Q}^+ \rightarrow \mathbb{R}$.
 $\frac{a}{b} \mapsto \log \max(|a|, |b|)$.

Will write down a local decomposition.
• First, given a prime p & $x \in \mathbb{Q}$, define

$|x|_p = p^{-n}$ where n maximal integer s.t. $x \cdot p^{-n}$ has no p in denominator

eg. $|12|_2 = \frac{1}{4}, |12|_3 = \frac{1}{3}, |12|_5 = 1,$

$|\frac{1}{12}|_2 = 4, |\frac{1}{12}|_3 = 3, |\frac{1}{12}|_5 = 1, \dots$

Set $|x|_\infty = \text{'usual abs. value'}$.

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Exercise: $h(x) = \sum_{p \in \{primes\} \cup \{\infty\}} \log \max(|x|_p, 1)$. (from $\mathbb{Q}^* \rightarrow \mathbb{R}$).

'local' = 'prime-by-prime'. We say h is an LD height.

[Formal def: After replacing x by alteration, h given by cdy metrized hermitian line bundle on a proper flat \mathbb{Z} -model.]

Thm [Néron, '65]: $\hat{h} : E(\mathbb{Q}) \rightarrow \mathbb{R}$ is an LD height.

With this notion of LDhts, let's go back to families:

Families II, Better

$$y^2 = x^3 - t^2x + t^3. \quad \bullet \text{ Fix } t \rightarrow \text{ell-curve } E_t$$

• Notice that (t, t) is always a rational pt on E_t ('section of the family').

Use this to define a fctn

$$H: \mathbb{Q}^* \longrightarrow \mathbb{R} \\ t \longmapsto \hat{h}_{E_t}(t, t).$$

$$\text{Eg. } 1 \longmapsto \hat{h}_E(p, 1), \quad E: y^2 = x^3 - x + 1, \quad p = (1, 1)$$

This is a ht by def'n. Is it an LDht? No.

Thm [Silverman, Tate]: There is an LD height $\hat{H}: \mathbb{Q}^* \rightarrow \mathbb{R}$

such that $H - \hat{H}$ is bounded.

(for any family, & any section).

Cor: $\{t \in \mathbb{Q} \mid (t, t) \text{ is torsion in } E_t(\mathbb{Q})\}$ is finite.

we gave a concept
w. R de Jong & O Brieskorn
using study of resistance
in electrical networks

Families with more parameters.

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Eg: $y^2 = x^3 - t^2x + s^2$, Fix $s, t \in \mathbb{Q}$ (st: $27s^4 - 4t^6 \neq 0$.)
 \rightarrow elliptic curve $E_{s,t}$.

"section" (t, s) .

In same way, get

$$H: \mathbb{Q} \times \mathbb{Q} \setminus \{27s^4 - 4t^6 = 0\} \rightarrow \mathbb{R}$$

$$(s, t) \longmapsto \hat{h}_{E_{s,t}}(t, s)$$

Again, $H(1,1) = \hat{h}_E(p)$, $E: y^2 = x^3 - x + 1$, $p = (1,1)$.

Noticed: If for all families "like this" (a Ab-sch / varieties @) there is an LD ht \hat{H} st. $H - \hat{H}$ bounded, then STC follows immediately.

Thm [H3]: The fctn H as above does NOT have bounded difference from any LD ht.

(Actually a complete classification of when this is possible. ~~Don't~~ Don't know an easy proof.)

Q: Is there an LD height \hat{H} such that $H - \hat{H}$ is "small"?

Dream: prove some cases of STC in this way.